CPE 323: The MSP430 Instruction Set Architecture

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Outline

- Introduction
- Registers
- Memory
- **Addressing Modes**
- Instruction Set
- Instruction Formats and Encodings



Memory

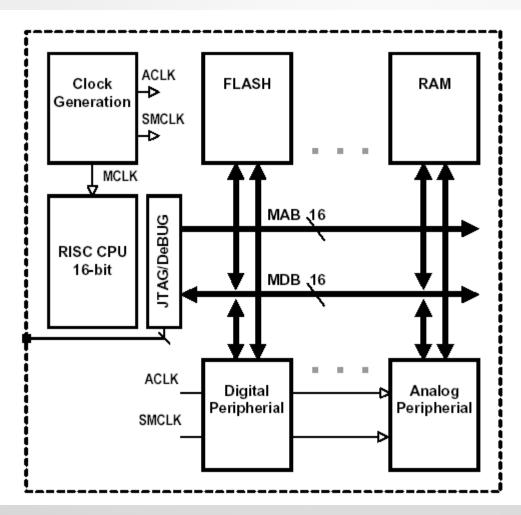
Intro



MSP 430 Modular Architecture

von-Neumann common bus connects CPU to all memory and peripherals

Embedded
emulation
accessed
in-application
with JTAG



Architecture reduces power consuming, noise generating fetches to memory

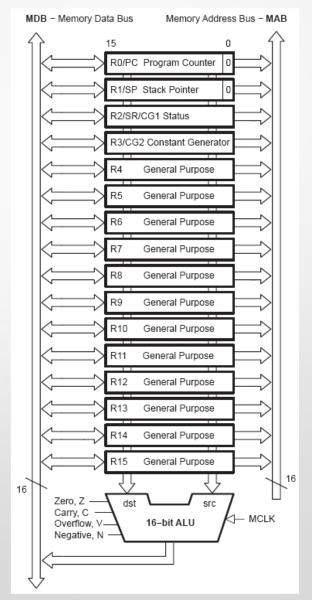
16-bit bus handles widewidth data much more effectively



Intro

MSP430 16-bit RISC

- Large 16-bit register file
- High-bandwidth 16-bit data and address bus
- RISC architecture with 27 instructions and 7 addressing modes
- Single-cycle register operations with full-access
- Direct memory-memory transfer designed for modern programming







Registers

15 0 0	RO - PC Program Counter 16-bit = no paging
15 0	R1 - SP Stack Pointer Addressable = great "C" code
reserved for future enhancements V SCG1 SCG0 Osc Off Off Off Off Off Off Off Off Off Of	R2 - SR Status Register Define LPMx
00000h 00001h 00002h 0004h 00008h 0FFFFh	R3/R2 - CG Constant Generator automatic generation of common used values reduces code size 30%
15 0	R4 - General Purpose
15 0	R15 - General Purpose

R4 through R15 are single-cycle, general purpose and identical in all respects - used for math, storage, and addressing modes.





PC/R0 – Program Counter



- The 16-bit program counter (PC/R0) points to the next instruction to be executed
- Each instruction uses an even number of bytes (two, four, or six), and the PC is incremented accordingly
 - Instruction accesses in the 64-KB address space are performed on word boundaries, and the PC is aligned to even addresses
- PC can be addressed by all instructions and all addressing modes
 - MOV #LABEL,PC; Branch to address LABEL
 - MOV LABEL,PC; Branch to address contained in LABEL
 - MOV @R14,PC; Branch indirect to address in R14

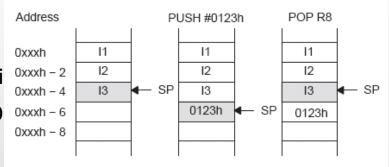




SP/R1 – Stack Pointer



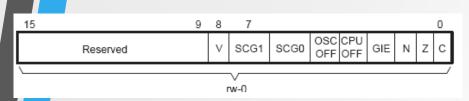
- The stack pointer (SP/R1) points to the current top of the stack
 - It uses a predecrement, postincrement scheme (SP points to the top of the stack, and the stack grows toward lower addresses)
- SP can be used by software with all instructions and addressing modes
- Examples
 - MOV 2 (SP), R6; Item I2 -> R6
 - MOV R7,0(SP); Overwrite TOS wi
 - PUSH #0123h; Put 0123h onto TO
 - POP R8; R8 = 0123h



• Question: Illustrate the stack contents after PUSH SP and POP SP instructions are executed?

SR/R2 – Status Register

Instruction Formats



 The status register (SR/R2), used as a source or destination register, can be used in the register mode only addressed with word instructions

<u>Addressing</u>

The remaining combinations of addressing modes are used to support the constant generator

Bit	Description	Description					
V	Overflow bit. This bit is set when overflows the signed-variable range	the result of an arithmetic operation le.					
	ADD(.B),ADDC(.B)	Set when: Positive + Positive = Negative Negative + Negative = Positive, otherwise reset					
	SUB(.B),SUBC(.B),CMP(.B)	Set when: Positive – Negative = Negative Negative – Positive = Positive, otherwise reset					
SCG1	System clock generator 1. This bit	, when set, turns off the SMCLK.					
SCG0	System clock generator 0. This bit, when set, turns off the DCO dc generator, if DCOCLK is not used for MCLK or SMCLK.						
OSCOFF	Oscillator Off. This bit, when set, turns off the LFXT1 crystal oscillator, when LFXT1CLK is not use for MCLK or SMCLK						
CPUOFF	CPU off. This bit, when set, turns of	off the CPU.					
GIE	General interrupt enable. This interrupts. When reset, all maskab	bit, when set, enables maskable le interrupts are disabled.					
N	Negative bit. This bit is set when the is negative and cleared when the i	ne result of a byte or word operation result is not negative.					
	Word operation:	N is set to the value of bit 15 of the result					
	Byte operation:	N is set to the value of bit 7 of the result					
Z	Zero bit. This bit is set when the re and cleared when the result is not	esult of a byte or word operation is 0 0.					
С	Carry bit. This bit is set when the produced a carry and cleared whe	e result of a byte or word operation n no carry occurred.					



Constant Generation

- Six commonly-used constants are generated with the constant generator registers R2 and R3
 - Adv.: No special instructions, no special code, no extra memory access
- Assembler uses the constant generator automatically if one of the six constants is used as an immediate source operand.
 Registers R2 and R3, used in the constant mode, cannot be addressed explicitly; they act as source-only registers.
- The constants are selected with the source-register addressing modes (As), as described below.

Register	As	Constant	Remarks
R2	00		Register mode
R2	01	(0)	Absolute address mode
R2	10	00004h	+4, bit processing
R2	11	00004H	
			+8, bit processing
R3	00	00000h	0, word processing
R3	01	00001h	+1
R3	10	00002h	+2, bit processing
R3	11	0FFFFh	–1, word processing



Constant Generation

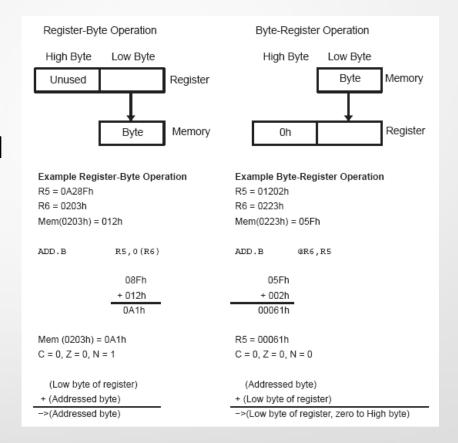
- Constant generator allows for additional 24 instructions that are emulated
- Examples
 - CLR dst MOV R3,dst
 - INC dst ADD 0(R3),dst



Memory

General-Purpose Registers

• The twelve registers, R4–R15, are general-purpose registers. All of these registers can be used as data registers, address pointers, or index values and can be accessed with byte or word instructions as shown below







Memory Organization

- Word alignment
 - Bytes are located at even or odd addresses
 - Words are only located at even addresses
- Endianess (little-endian)
 - When using word instructions, only even addresses may be used
 - The Low byte of a word is always an even address
 - The high byte is at the next odd address
 - For example, if a data word is located at address xxx4h, then the low byte of that data word is located at address xxx4h, and the high byte of that word is located at address xxx5h

		•••			xxxAh
15	14	Bits	9	8	xxx9h
7	6	Bits	1	0	xxx8h
		Byte			xxx7h
		Byte			xxx6h
	W	ord (High By	te)		xxx5h
	W	ord (Low By	te)		xxx4h
	xxx3h				

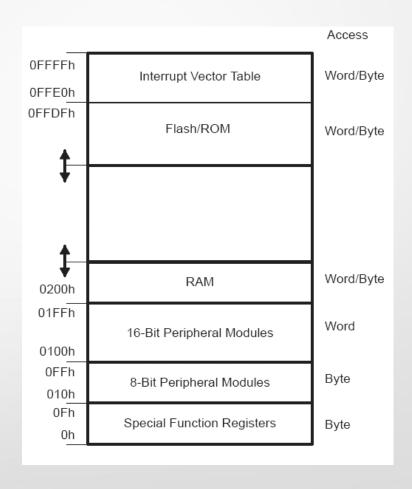


Intro



Address Space

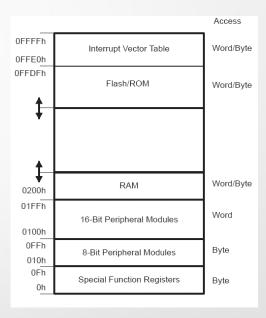
- von-Neumann architecture: one address space shared with special function registers (SFRs), peripherals, RAM, and Flash/ROM memory as shown
- Memory maps are device specific
- Code access are always performed on even addresses
- Data can be accessed as bytes or words
- The addressable memory space is 64 KB
 - Extended architecture allows for 1 MB address space





Address Space (cont'd)

- Special Function Registers (SFRs)
 - Some peripheral functions are configured in the SFRs
 - The SFRs are located in the lower 16 bytes of the address space, and are organized by byte
 - SFRs must be accessed using byte instructions only
- Peripheral modules (PM)
 - Peripheral modules are mapped into the address space
 - Address space 0100-01FFh is reserved for 16-bit PMs
 - Should be accessed with word instructions
 - If byte instructions are used, only even addresses are permissible, and the high byte of the result is always 0
 - Address space 010h-0FFh is reserved for 8-bit PMs
 - Should be accessed with byte instructions
 - Read access of byte modules using word instructions results in unpredictable data in the high byte
 - If word data is written to a byte module only the low byte is written into the peripheral register, ignoring the high byte

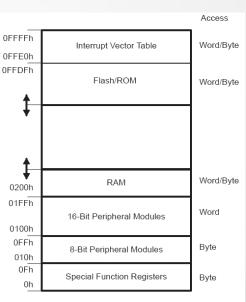






Address Space (cont'd)

- RAM
 - RAM starts at 0200h
 - End address of RAM depends on the amount of RAM present and varies by device
 - RAM can be used for both code and data
- Flash/ROM
 - Start address of Flash/ROM depends on the amount of Flash/ROM present and varies by device
 - End address for Flash/ROM is 0FFFFh
 - Flash can be used for both code and data;
 Word or byte tables can be stored and used in Flash/ROM without the need to copy the tables to RAM before using them
- Interrupt vector table
 - Is mapped into the upper 16/32 words of Flash/ROM address space, with the highest priority interrupt vector at the highest Flash/ROM word address (OFFFEh)







Addressing Modes

- Register (a.k.a., register direct)
- Indexed
- Symbolic (a.k.a., PC relative, special case of indexed)
- Absolute (special case of indexed)
- Register indirect
- Register indirect with autoincrement
- Immediate

Addressing Modes

 Seven addressing modes for the source operand and four addressing modes for the destination operand can address the complete address space with no exceptions.

Table 2. Address Mode Descriptions

ADDRESS MODE	s	D	SYNTAX	EXAMPLE	OPERATION
Register	● ● MOV Rs,Rd		MOV Rs,Rd	MOV R10,R11	R10> R11
Indexed	•	•	MOV X(Rn),Y(Rm)	MOV 2(R5),6(R6)	M(2+R5)> M(6+R6)
Symbolic (PC relative)	•	•	MOV EDE, TONI		M(EDE)> M(TONI)
Absolute	te ● MOV &MEM,&TCDAT		MOV &MEM,&TCDAT		M(MEM)> M(TCDAT)
Indirect	•		MOV @Rn,Y(Rm)	MOV @R10,Tab(R6)	M(R10)> M(Tab+R6)
Indirect autoincrement	MOV @Rn+ Rm		MOV @R10+,R11	M(R10)> R11 R10 + 2> R10	
Immediate	•		MOV #X,TONI	MOV #45,TONI	#45> M(TONI)

NOTE: S = source D = destination





Addressing Modes

Instruction Formats

The bit numbers in the table below describe the contents of the As (source) and Ad (destination) mode bits

As/Ad	Addressing Mode	Syntax	Description
00/0	Register mode	Rn	Register contents are operand
01/1	Indexed mode	X(Rn)	(Rn + X) points to the operand. X is stored in the next word.
01/1	Symbolic mode	ADDR	(PC + X) points to the operand. X is stored in the next word. Indexed mode X(PC) is used.
01/1	Absolute mode	&ADDR	The word following the instruction contains the absolute address. X is stored in the next word. Indexed mode X(SR) is used.
10/-	Indirect register mode	@Rn	Rn is used as a pointer to the operand.
11/-	Indirect autoincrement	@Rn+	Rn is used as a pointer to the operand. Rn is incremented afterwards by 1 for .B instructions and by 2 for .W instructions.
11/-	Immediate mode	#N	The word following the instruction contains the immediate constant N. Indirect autoincrement mode @PC+ is used.





Register Addressing Mode

Op-Code	Source-Regist	er Ad	B/W	As	Destination-Register
0100	0100	0	0	00	0101
4405	mov.w	R4,R5			;
4445	mov.b	R4,R5			;

Valid for Source and destination As=00, Ad=0

The operand is contained in one of the CPU registers R0 to R15. This is the fastest addressing mode and needs the least memory .



Register Addressing Mode (cont'd)

Example: MOV R10,R11 Before: After: 0A023h R10 0A023h R10 0FA15h R11 R11 0A023h PC_{old} PC $PC_{old} + 2$ PC

Note: Data in Registers

The data in the register can be accessed using word or byte instructions. If byte instructions are used, the high byte is always 0 in the result. The status bits are handled according to the result of the byte instruction.



Register-Indexed Addressing Mode

	Op-Code	Source-Register	r Ad	B/W	As	Destination-Register
•	0100	0100	1	0	01	0101
	4495 01000200	mov.w 1	00h(R	24),	200h (R	5) ;

44150100 mov.w 100h(R4),R5 ;

Valid for Source and destination As=01, Ad=1

The address of the operand is the sum of the index and the contents of the register.



Register-Indexed Addressing Mode (cont'd)

Example: MOV 2(R5),6(R6);								
Before:	Address Space	Register	After:	Address Space	Register			
0FF16h 0FF14h 0FF12h	00006h 00002h 04596h	R5 01080h R6 0108Ch PC	0FF16h 0FF14h 0FF12h	0xxxxh 00006h 00002h 04596h	R5 01080h R6 0108Ch			
01094h 01092h 01090h	Oxxxxh O5555h Oxxxxh	0108Ch +0006h 01092h	01094h 01092h 01090h	0xxxxh 01234h 0xxxxh				
01084h 01082h 01080h	0xxxxh 01234h 0xxxxh	01080h +0002h 01082h	01084h 01082h 01080h	Oxxxxh O1234h Oxxxxh				





Symbolic Addressing Mode

Instruction Formats

Op-Code	Source-Register		Ad	B/W	As	Destination-Register
0100	0000		1	0	01	0000
4<u>0</u>9<u>0</u> ffa80006	mov.w	EDE,	ТО	NI		;
4015ffac	mov.w	EDE,	R5			;

Source and destination As=01, Ad=1

Addressing

The content of the addresses EDE / TONI are used for the operation. The source or destination address is computed as a difference from the PC and uses the PC in indexed addressing mode. Any address in the 64k memory space is addressable.



Symbolic Addressing Mode (cont'd)

Example:	MOV EDE, TONI ; Source address EDE = 0F016h ; Dest. address TONI=01114h							
Before:	Address Space	Registe		After:	Address Space	Register I PC		
0FF16h	011FEh			0FF16h	011FEh			
0FF14h	0F102h			0FF14h	0F102h			
0FF12h	04090h	PC		0FF12h	04090h			
0F018h 0F016h 0F014h	0xxxxh 0A123h 0xxxxh	0FF +0F10 0F0	02h	0F018h 0F016h 0F014h	0xxxxh 0A123h 0xxxxh			
01116h 01114h 01112h	0xxxxh 05555h 0xxxxh	0FF +011F 011	Eh_	01116h 01114h 01112h	0xxxxh 0A123h 0xxxxh			





Absolute Addressing Mode

Instruction Formats

Op-Code	Source-Regis	ster	Ad	B/W	As	Destination-Register
0100	0010		1	0	01	0010
4 <u>2</u> 9 <u>2</u> 01720174	mov.w	&CC	RO,	&CC	R1	;
42150172	mov.w	&CC	RO,	R5		;

Source and destination As=01, Ad=1

The contents of the fixed addresses are used for the operation. The SR is used in the indexed mode to create an absolute 0. Use for hardware peripherals located at an absolute address that can never be relocated.





Absolute Addressing Mode (cont'd)

Example:	MOV	&EDE,&TONI	;Source ad;		
Before:	Address Space	Register	After:	Address Space	Register
0FF16h	01114h		0FF16h	0xxxxh 01114h	PC
0FF14h	0F016h		0FF14h	0F016h	
0FF12h	04292h	PC	0FF12h	04292h	
				\square	
0F018h	0xxxxh		0F018h	0xxxxh	
0F016h	0A123h		0F016h	0A123h	
0F014h	0xxxxh		0F014h	0xxxxh	
					1
01116h	0xxxxh		01116h	0xxxxh	
01114h	01234h		01114h	0A123h	
01112h	0xxxxh		01112h	0xxxxh	



Register Indirect Addressing Mode

Instruction Formats

Op-Code	Source-Register	Ad	B/W	As	Destination-Register
0100	0100	0	0	10	0101
4425	mov.w @R	4,R5			;
4465	mov.b @R	4,R5			;

Source only As=10, Ad=n/a

The registers are used as a pointer to the operand.

The indexed mode with zero index may be used for "indirect register addressing" of the destination operand.

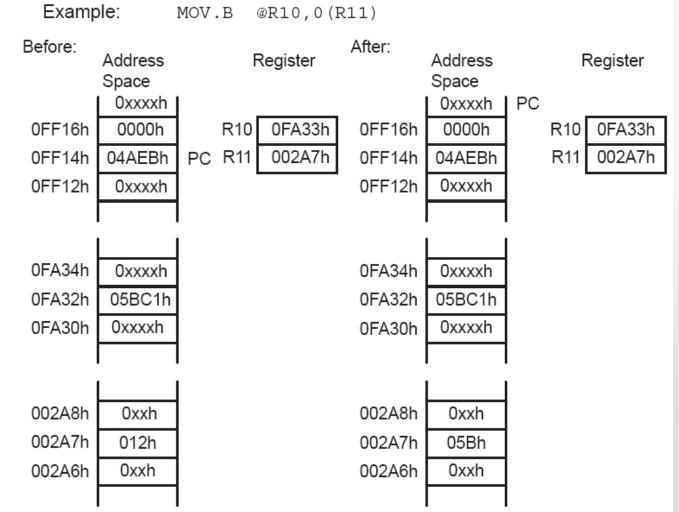
44a50000

mov.w

@R4,0(R5)



Register Indirect Addressing Mode (cont'd)





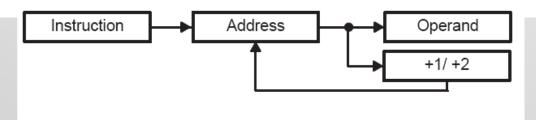
Register Indirect Autoincrement AM

Instruction Formats

Op-Code	Source-Register	Ad	B/W	As	Destination-Register
0100	0100	0	0	11	0101
4435	mov.w @F	4+,R	5		;
4475	mov.b @R	4+,R	5		;

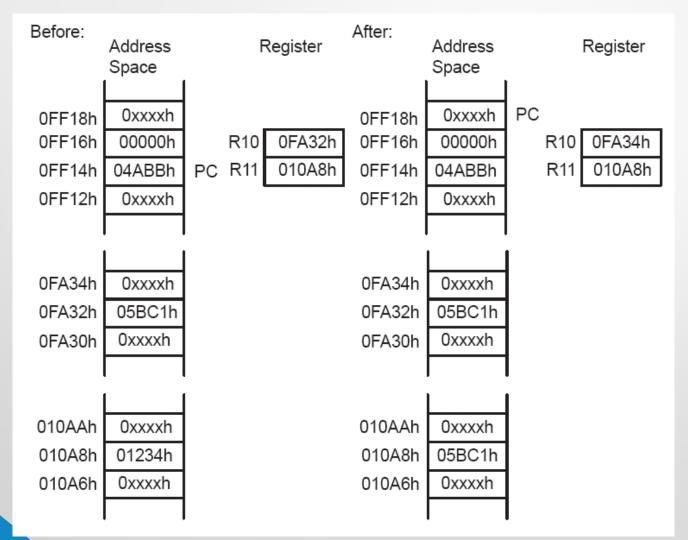
Source only As=11, Ad=n/a

The registers are used as a pointer to the operand. The registers are incremented afterwards - by 1 in byte mode, by 2 in word mode.





Register Indirect Autoincrement AM (cont'd)





Memory

Immediate Addressing Mode

	Op-Code	Source-Register	Ad	B/W	As	Destination-Register
•	0100	0000	0	0	11	0101

4<u>0</u>351234

mov.w

#1234h,R5

; Any 16-bit value

Source only As=11, Ad=n/a

Any immediate 8 or 16 bit constant can be used with the instruction. The PC is used in autoincrement mode to emulate this addressing mode.



Immediate Addressing Mode (cont'd)

Example:	VOM	#45h,TONI				
Before:	Address Space	Register	After:	Address Space		Register
			0FF18h	0xxxxh	PC	
0FF16h	01192h		0FF16h	01192h		
0FF14h	00045h		0FF14h	00045h		
0FF12h	040B0h	PC	0FF12h	040B0h		
		0FF16h			1	
010AAh	0xxxxh	+01192h	010AAh	0xxxxh		
010A8h	01234h	010A8h	010A8h	00045h		
010A6h	0xxxxh		010A6h	0xxxxh		

Instruction Set

- 27 core instructions
 - Have unique op-codes decoded by the CPU
- 24 emulated instructions
 - Make code easier to write and read, but do not have op-codes; instead an equivalent core instruction is generated
 - No code or performance penalty for using emulated instructions
- 3 core instruction formats
 - **Dual-operand**
 - Single-operand
 - Jump
- All single- and dual-operand instructions can be byte or word instructions by using .B or .W (default) extensions
 - Byte instructions are used to access byte data or byte peripherals
 - Word instructions are used to access word data or word peripherals





27 Core RISC Instructions

Format I	Format II	Format III
Source, Destination	Single Operand	+/- 9bit Offset
add(.b)	call	jmp
addc(.b)	swpb	jc
and(.b)	sxt	jnc
bic(.b)	push(.b)	jeq
bis(.b)	reti	jne
bit(.b)	rra(.b)	jge
cmp(.b)	rrc(.b)	jl
dadd(.b)		jn
mov(.b)		
sub(.b)		
subc(.b)		
xor(.b)		





Emulated Instructions: Examples

```
// clear carry
clrc
                   ; Clear Carry bit
bic.w \#0x01, SR ; Clear Bit 0 in SR
// decrement
dec.w R14
          ; Decrement R14
sub.w \#0x01, R4 ; Core instruction
// return from subroutine
ret
mov.w @SP+, PC
```





51 Total Instructions

Format I	Format II	Format III	Support
Source, Destination	Single Operand	+/- 9bit Offset	
add(.b)	br	jmp	clrc
addc(.b)	call	jc	setc
and(.b)	swpb	jnc	clrz
bic(.b)	sxt	jeq	setz
bis(.b)	push(.b)	jne	clrn
bit(.b)	pop.(b)	jge	setn
cmp(.b)	rra(.b)	jl	dint
dadd(.b)	rrc(.b)	jn	eint
mov(.b)	inv(.b)		nop
sub(.b)	inc(.b)		ret
subc(.b)	incd(.b)		reti
xor(.b)	dec(.b)		
	decd(.b)		
	adc(b)		
	sbc(.b)		
	clr(.b)		
	dadc(.b)		
	rla(.b)		
	rlc(.b)		
	tst(.b)		



gisters Memory Addressing Instruction Set Instruction Format



Double operand instructions

15 14 13	12	11	10	9	8	7	6	5	4	3	2	1	0
Op-code			S-R	eg		Ad	B/W		As		D-R	eg	

Mnemonic	S-Reg,	Operation		State	us Bit	s
	D-Reg		٧	N	Z	С
MOV(.B)	src,dst	$src \to dst$	_	_	_	_
ADD(.B)	src,dst	$\text{src} + \text{dst} \rightarrow \text{dst}$	*	*	*	*
ADDC(.B)	src,dst	$src + dst + C \rightarrow dst$	*	*	*	*
SUB(.B)	src,dst	$dst + .not.src + 1 \rightarrow dst$	*	*	*	*
SUBC(.B)	src,dst	$\text{dst + .not.src + C} \rightarrow \text{dst}$	*	*	*	*
CMP(.B)	src,dst	dst – src	*	*	*	*
DADD(.B)	src,dst	$\text{src + dst + C} \rightarrow \text{dst (decimally)}$	*	*	*	*
BIT(.B)	src,dst	src .and. dst	0	*	*	*
BIC(.B)	src,dst	.not.src .and. dst \rightarrow dst	_	_	_	_
BIS(.B)	src,dst	$\text{src .or. dst} \rightarrow \text{dst}$	_	_	_	_
XOR(.B)	src,dst	$\text{src .xor. dst} \to \text{dst}$	*	*	*	*
AND(.B)	src,dst	src .and. dst $ ightarrow$ dst	0	*	*	*



Single Operand Instruction

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			0	p-code	Э				B/W		Ad		D/S-	Reg	

Mnemonic	S-Reg,	Operation	Stat	us Bi	ts	
	D-Reg		V	N	Z	С
RRC(.B)	dst	$C \to MSB \to \dots \dots LSB \to C$	*	*	*	*
RRA(.B)	dst	$MSB \to MSB \to LSB \to C$	0	*	*	*
PUSH(.B)	src	$SP-2 \rightarrow SP,src \rightarrow @SP$	-	_	_	_
SWPB	dst	Swap bytes	-	_	_	_
CALL	dst	$SP-2 \rightarrow SP, PC\text{+}2 \rightarrow @SP$	-	_	_	_
		$\text{dst} \to \text{PC}$				
RETI		$TOS \rightarrow SR,SP + 2 \rightarrow SP$	*	*	*	*
		$TOS \rightarrow PC, SP + 2 \rightarrow SP$				
SXT	dst	Bit $7 \rightarrow$ Bit 8Bit 15	0	*	*	*





Jump Instructions

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Ор-сос	le		С						10-Bi	t PC C	Offset			

Mnemonic	S-Reg, D-Reg	Operation
JEQ/JZ	Label	Jump to label if zero bit is set
JNE/JNZ	Label	Jump to label if zero bit is reset
JC	Label	Jump to label if carry bit is set
JNC	Label	Jump to label if carry bit is reset
JN	Label	Jump to label if negative bit is set
JGE	Label	Jump to label if (N .XOR. V) = 0
JL	Label	Jump to label if (N .XOR. V) = 1
JMP	Label	Jump to label unconditionally





3 Instruction Formats

;	Format	Ι	Source	and	Destination
---	--------	---	--------	-----	-------------

	Op-Code	Source-Register		Source-Register Ad B/W As		As	Desti	nation-Register	
5	405	add.w	R4,F	15			;	R4+R5=R5	XXXX
<u></u>	3445	add.b	R4,R	25			;	R4+R5=R5	0.0xx

; Format II Destination Only

	Op-Code	B/W	Ad	D/S- Register	
6404	rlc.w	R4			;
6444	rlc.b	R4			;

Format III There are 8(Un)conditional Jumps

	Op-Code	Condition	10-bit PC offs	set	
3	3c28	jmp	Loop_1	;	Goto Loop_

$$PC_{new} = PC_{old} + 2 + PC_{offset} \times 2$$



Instruction Cycles and Lengths

Instruction Formats

- The number of CPU clock cycles required for an instruction depends on the instruction format and the addressing modes used - not the instruction itself
- The number of clock cycles refers to the MCLK



Registers Memory Addressing Instruction Set Instruction Formats

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Format I: Instruction Cycles and Length

Table 3–16. Format 1 Instruction Cycles and Lengths

Addres	sing Mode	No. of	Length of		
Src	Dst	Cycles	Instruction		Example
Rn	Rm	1	1	MOV	R5,R8
	PC	2	1	BR	R9
	x(Rm)	4	2	ADD	R5,4(R6)
	EDE	4	2	XOR	R8,EDE
	&EDE	4	2	MOV	R5, &EDE
@Rn	Rm	2	1	AND	@R4,R5
	PC	2	1	BR	@R8
	x(Rm)	5	2	XOR	@R5,8(R6)
	EDE	5	2	MOV	@R5,EDE
	&EDE	5	2	XOR	@R5,&EDE
@Rn+	Rm	2	1	ADD	@R5+,R6
	PC	3	1	BR	@R9+
	x(Rm)	5	2	XOR	@R5,8(R6)
	EDE	5	2	MOV	@R9+,EDE
	&EDE	5	2	MOV	@R9+,&EDE
#N	Rm	2	2	MOV	#20,R9
	PC	3	2	BR	#2AEh
	x(Rm)	5	3	MOV	#0300h,0(SP)
	EDE	5	3	ADD	#33,EDE
	&EDE	5	3	ADD	#33,&EDE
x(Rn)	Rm	3	2	MOV	2(R5),R7
	PC	3	2	BR	2 (R6)
	TONI	6	3	MOV	4(R7),TONI
	x(Rm)	6	3	ADD	4(R4),6(R9)
	&TONI	6	3	MOV	2(R4),&TONI
EDE	Rm	3	2	AND	EDE,R6
	PC	3	2	BR	EDE
	TONI	6	3	CMP	EDE, TONI
	x(Rm)	6	3	MOV	EDE,0(SP)
	&TONI	6	3	MOV	EDE, &TONI
&EDE	Rm	3	2	MOV	&EDE,R8
	PC	3	2	BRA	&EDE
	TONI	6	3	MOV	&EDE, TONI
	x(Rm)	6	3	MOV	&EDE,0(SP)
	&TONI	6	3	MOV	&EDE, &TONI



Format II and Format III: Instruction Cycles and Length

Table 3-15. Format-II Instruction Cycles and Lengths

	No.	of Cycles		_	
Addressing Mode	RRA, RRC SWPB, SXT	PUSH	CALL	Length of Instruction	Example
Rn	1	3	4	1	SWPB R5
@Rn	3	4	4	1	RRC @R9
@Rn+	3	5	5	1	SWPB @R10+
#N	(See note)	4	5	2	CALL #0F000h
X(Rn)	4	5	5	2	CALL 2(R7)
EDE	4	5	5	2	PUSH EDE
&EDE	4	5	5	2	SXT &EDE

- Format III: all jump instructions take 2 clock cycles to execute and are 1 word long
- Interrupt and reset cycles

Table 3-14. Interrupt and Reset Cycles

Action	No. of Cycles	Length of Instruction
Return from interrupt (RETI)	5	1
Interrupt accepted	6	-
WDT reset	4	-
Reset (RST/NMI)	4	-





Instruction Encoding

	000	040	080	0C0	100	140	180	1C0	200	240	280	2C0	300	340	380	3C0
0xxx																
4xxx																
8xxx																
Cxxx																
1xxx	RRC	RRC.B	SWPB		RRA	RRA.B	SXT		PUSH	PUSH.B	CALL		RETI			
14xx																
18xx																
1Cxx																
20xx	JNE/JNZ															
24xx	JEQ/JZ															
28xx	JNC															
2Cxx	JC															
30xx	JN															
34xx							J	IGE								
38xx								IL								
3Схх							J	MP								
4xxx							Ν	ИOV, N	IOV.B							
5xxx								ADD, A								
6xxx							A	ADDC,	ADDC.	В						
7xxx							5	SUBC,	SUBC.	В						
8xxx							9	SUB, S	UB.B							
9ххх							(CMP, C	MP.B							
Axxx								DADD,	DADD.	В						
Bxxx							Е	BIT, BIT	.В							
Cxxx							Е	BIC, BIG	C.B							
Dxxx							Е	BIS, BIS	S.B							
Exxx								(OR, X								
Fxxx							1	AND, A	ND B							





MSP 430 System Architecture: A Closer Look

