



CPE 323:

The MSP430 Assembly

Language Programming

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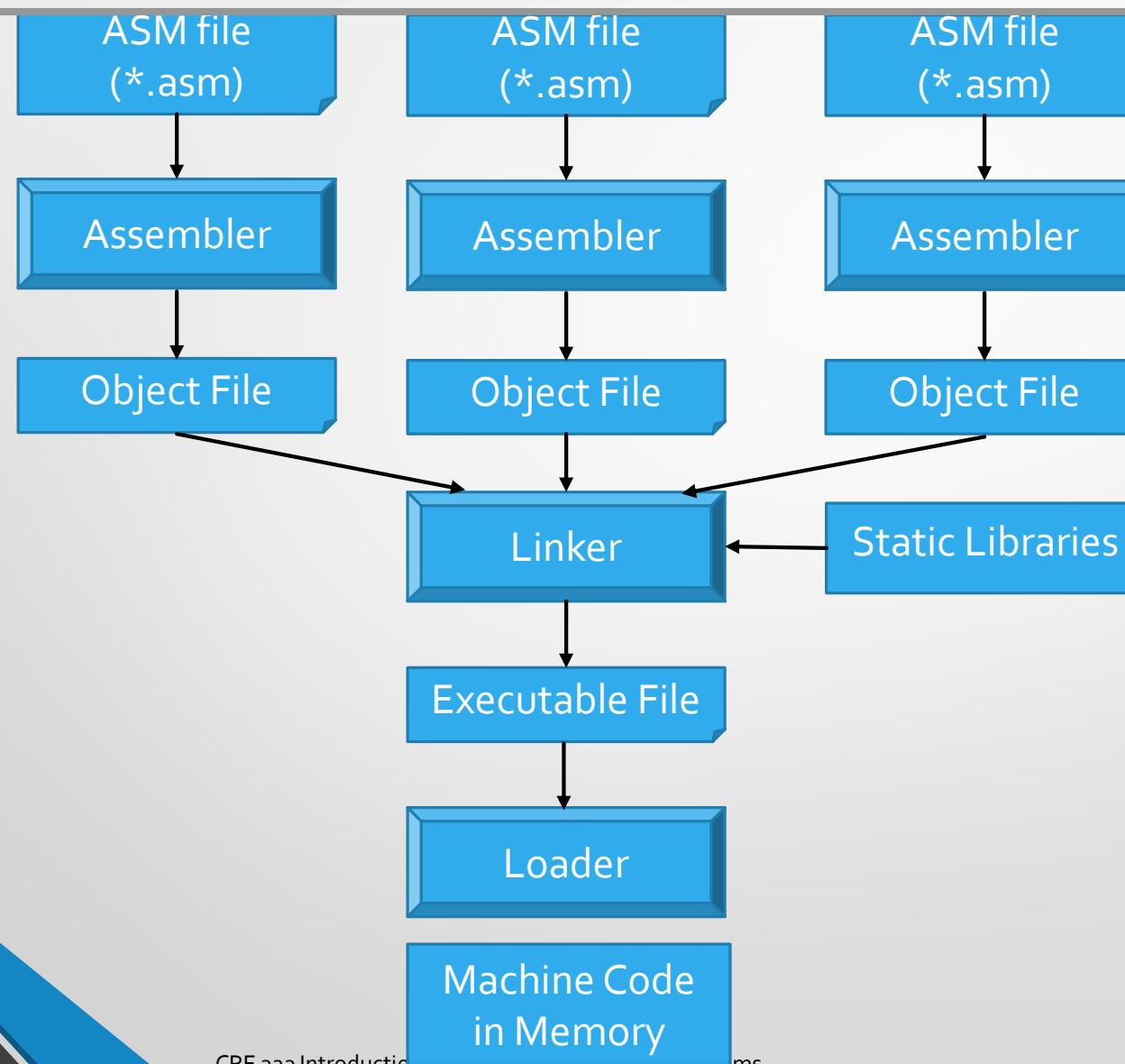
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Outline

- Introduction
- Assembly language directives
- SUMI/SUMD
 - Adding two 32-bit numbers (decimal, integers)
- CountEs: Counting characters ‘E’
- Subroutines
 - CALL&RETURN
 - Subroutine Nesting
 - Passing parameters
 - Stack and Local Variables
- Performance

Assembly Programming Flow



Assembly Directives

- Assembly language directives tell the assembler to
 - Set the data and program at particular addresses in address space
 - Allocate space for constants and variables
 - Define synonyms
 - Include additional files
 - ...
- Typical directives
 - Equate: assign a value to a symbol
 - Origin: set the current location pointer
 - Define space: allocate space in memory
 - Define constant: allocate space for and initialize constants
 - Include: loads another source file

ASM430 Section Control Directives

- CCStudio ASM430 has three predefined sections into which various parts of a program are assembled
 - .bss: Uninitialized data section
 - .data: Initialized data section
 - .text: Executable code section

Description	ASM430 (CCS)	A430 (IAR)
Reserve size bytes in the uninitialized sect.	.bss	-
Assemble into the initialized data section	.data	RSEG const
Assemble into a named initialized data sect.	.sect	RSEG
Assemble into the executable code	.text	RSEG code
Reserve space in a named (uninitialized) section	.usect	-
Align on byte boundary	.align 1	-
Align on word boundary	.align 2	EVEN

Examples

```
; IAR

    RSEG DAT16_N      ; switch to DATA segment
    EVEN                 ; make sure it starts at even address
MyWord: DS 2            ; allocate 2 bytes / 1 word
MyByte: DS 1            ; allocate 1 byte

; CCS Assembler (Example #1)
MyWord: .usect ".bss", 2, 2 ; allocate 1 word
MyByte: .usect ".bss", 1      ; allocate 1 byte

; CCS Assembler (Example #2)
.bss MyWord,2,2 ; allocate 1 word
.bss MyByte,1    ; allocate 1 byte
```

Constant Initialization Directives

Description	ASM430 (CCS)	A430 (IAR)
Initialize one or more successive bytes or text strings	.byte or .string	DB
Initialize 32-bit IEEE floating-point	.float	DF
Initialize a variable-length field	.field	-
Reserve size bytes in the current location	.space	DS
Initialize one or more 16-bit integers	.word	DW
Initialize one or more 32-bit integers	.long	DL

Directives: Dealing with Constants

```
b1:      .byte  5           ; allocates a byte in memory and initialize it with 5
b2:      .byte -122         ; allocates a byte with constant -122
b3:      .byte 10110111b    ; binary value of a constant
b4:      .byte 0xA0          ; hexadecimal value of a constant
b5:      .byte 123q          ; octal value of a constant
tf:      .equ   25
```

Word view of Memory

Label	Address	Memory[15:8]	Memory[7:0]
b1	0x3100	0x86	0x05
b3	0x3102	0xA0	0xB7
b5	0x3104	--	0x53

Byte view of Memory

Label	Address	Memory[7:0]
b1	0x3100	0x05
b2	0x3101	0x86
b3	0x3102	0xB7
b4	0x3103	0xA0
b5	0x3104	0x53

Directives: Dealing with Constants

```
...  
w1:      .word   21           ; allocates a word constant in memory;  
  
w2:      .word   -21  
w3:      .word tf  
dw1:     .long   100000        ; allocates a long word size constant in memory;  
                  ; 100000 (0x0001_86A0)  
dw2:     .long  0xFFFFFEA
```

Label	Address	Memory[15:8]	Memory[7:0]
w1	0x3106	0x00	0x15
w2	0x3108	0xFF	0xEB
w3	0x310A	0x00	0x19
dw1	0x310C	0x86	0xA0
	0x310E	0x00	0x01
dw2	0x3110	0xFF	0xEA
	0x3112	0xFF	0xFF

Directives: Dealing with Constants

```
s1:     .byte 'A', 'B', 'C', 'D' ; allocates 4 bytes in memory with string ABCD
s2:     .byte "ABCD", ' ' ; allocates 5 bytes in memory with string ABCD + NULL
```

Label	Address	Memory[15:8]	Memory[7:0]
s1	0x3114	0x42	0x41
	0x3116	0x44	0x43
s2	0x3118	0x42	0x41
	0x311A	0x44	0x43
	0x311C	--	0x00
	0x311E		

Table of Symbols

- Created by the assembler
(think about this as a table of synonyms)

Symbol	Value [hex]
b1	0x3100
b2	0x3101
b3	0x3102
b4	0x3103
b5	0x3104
tf	0x0019
w1	0x3106
w2	0x3108
w3	0x310A
dw1	0x310C
dw2	0x3110
s1	0x3114
s2	0x3118

Directives: Variables in RAM

```

.bss v1b,1,1      ; allocates a byte in memory, equivalent to DS 1
.bss v2b,1,1      ; allocates a byte in memory
.bss v3w,2,2      ; allocates a word of 2 bytes in memory
.bss v4b,8,2      ; allocates a buffer of 2 long words (8 bytes)
.bss vx,1,1
  
```

Label	Address	Memory[15:8]	Memory[7:0]
v1b	0x1100	--	--
v3w	0x1102	--	--
v4b	0x1104	--	--
	0x1106	--	--
	0x1108	--	--
	0x110A	--	--
vx	0x110C		

Symbol	Value [hex]
v1b	0x1100
v2b	0x1101
v3w	0x1102
v4b	0x1104
vx	0x110C

Decimal/Integer Addition of 32-bit Numbers

- Problem
 - Write an assembly program that finds a sum of two 32-bit numbers
 - Input numbers are decimal numbers (8-digit in length)
 - Input numbers are signed integers in two's complement
- Data:
 - `lint1: DC32 0x45678923`
 - `lint2: DC32 0x23456789`
 - Decimal sum: `0x69135712`
 - Integer sum: `0x68ac31ac`
- Approach
 - Input numbers: storage, placement in memory
 - Results: storage (ABSOLUTE ASSEMBLER)
 - Main program: initialization, program loops
 - Decimal addition, integer addition

Decimal/Integer Addition of 32-bit Numbers

```
;-----  
; File      : LongIntAddition.asm  
; Function   : Sums up two long integers represented in binary and BCD  
; Description: Program demonstrates addition of two operands lint1 and lint2.  
;  
;           Operands are first interpreted as 32-bit decimal numbers and  
;           and their sum is stored into lsumd;  
;  
;           Next, the operands are interpreted as 32-bit signed integers  
;           in two's complement and their sum is stored into lsumi.  
;  
; Input      : Input integers are lint1 and lint2 (constants in flash)  
; Output     : Results are stored in lsumd (decimal sum) and lsumi (int sum)  
; Author     : A. Milenkovic, milenkovic@computer.org  
; Date       : August 24, 2018  
;  
        .cdecls C,LIST,"msp430.h"          ; Include device header file  
  
;  
        .def    RESET                   ; Export program entry-point to  
;                                     ; make it known to linker.  
;  
        .text                          ; Assemble into program memory.  
        .retain                         ; Override ELF conditional linking  
;                                     ; and retain current section.  
        .retainrefs                     ; And retain any sections that have  
;                                     ; references to current section.  
;
```

Decimal/Integer Addition of 32-bit Numbers (cont'd)

```
lint1:.long 0x45678923
lint2:.long 0x23456789
;
;
lsumd:.usect ".bss", 4,2 ; allocate 4 bytes for decimal result
lsumi:.usect ".bss", 4,2 ; allocate 4 bytes for integer result
;
RESET:    mov.w   #__STACK_END,SP           ; Initialize stack pointer
StopWDT:  mov.w   #WDTPW|WDTHOLD,&WDTCTL ; Stop watchdog timer
;
```

Decimal/Integer Addition of 32-bit Numbers (cont'd)

```
;-----  
; Main code here  
;  
        clr.w   R2          ; clear status register  
        mov.w   lint1, R8    ; get lower 16 bits from lint1 to R8  
        dadd.w  lint2, R8    ; decimal addition, R8 + lower 16-bit of lint2  
        mov.w   R8, lsumd    ; store the result (lower 16-bit)  
        mov.w   lint1+2, R8   ; get upper 16 bits of lint1 to R8  
        dadd.w  lint2+2, R8   ; decimal addition  
        mov.w   R8, lsumd+2  ; store the result (upper 16-bit)  
        mov.w   lint1, R8    ; get lower 16 bite from lint1 to R8  
        add.w   lint2, R8    ; integer addition  
        mov.w   R8, lsumi    ; store the result (lower 16 bits)  
        mov.w   lint1+2, R8   ; get upper 16 bits from lint1 to R8  
        addc.w  lint2+2, R8   ; add upper words, plus carry  
        mov.w   R8, lsumi+2  ; store upper 16 bits of the result  
  
        jmp $                ; jump to current location '$'  
                           ; (endless loop)
```

Decimal/Integer Addition of 32-bit Numbers (cont'd)

```
;-----  
; Stack Pointer definition  
;  
        .global __STACK_END  
        .sect    .stack  
  
;  
; Interrupt Vectors  
;  
        .sect    ".reset"           ; MSP430 RESET Vector  
        .short   RESET
```

Version 2: Decimal/Integer Addition of 32-bit Numbers (cont'd)

```
; Decimal addition

        mov.w  #lint1, R4          ; pointer to lint1
        mov.w  #lsumd, R8          ; pointer to lsumd
        mov.w  #2, R5              ; R5 is a counter (32-bit=2x16-bit)
        clr.w  R10                ; clear R10
lda:   mov.w  4(R4), R7          ; load lint2
        mov.w  R10, R2              ; load original SR
        dadd.w @R4+, R7            ; decimal add lint1 (with carry)
        mov.w  R2, R10              ; backup R2 in R10
        mov.w  R7, 0(R8)            ; store result (@R8+0)
        add.w  #2, R8                ; update R8
        dec.w  R5                  ; decrement R5
        jnz   lda                  ; jump if not zero to lda
```

Version 2: Decimal/Integer Addition of 32-bit Numbers (cont'd)

```
; Integer addition

        mov.w  #lint1, R4          ; pointer to lint1
        mov.w  #lsumi, R8          ; pointer to lsumi
        mov.w  #2, R5              ; R5 is a counter (32-bit=2x16-bit)
        clr.w  R10                ; clear R10
lia:     mov.w  4(R4), R7          ; load lint2
        mov.w  R10, R2              ; load original SR
        addc.w @R4+, R7            ; decimal add lint1 (with carry)
        mov.w  R2, R10              ; backup R2 in R10
        mov.w  R7, 0(R8)            ; store result (@R8+0)
        add.w  #2, R8                ; update R8
        dec.w  R5                  ; decrement R5
        jnz   lia                  ; jump if not zero to lia

        jmp   $                      ; jump to current location '$'
                                    ; (endless loop)
```

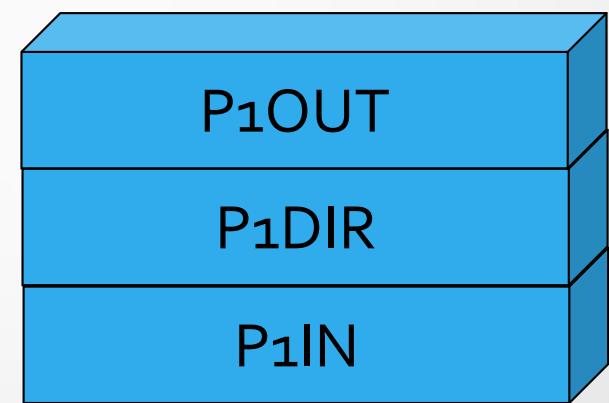
Count Characters ‘E’

- Problem
 - Write an assembly program that processes an input string to find the number of characters ‘E’ in the string
 - The number of characters is “displayed” on the port 1 of the MSP430
- Example:
 - mystr=“HELLO WORLD, I AM THE MSP430!”, ”
 - P1OUT=0x02
- Approach
 - Input string: storage, placement in memory
 - Main program: initialization, main program loop
 - Program loop: iterations, counter, loop exit
 - Output: control of ports

Programmer's View of Parallel Ports

- Parallel ports: $x=1,2,3,4,5, \dots$
- Each can be configured as:
 - Input: $PxDIR=0x00$ (default)
 - Output: $PxDIR=0xFF$
- Writing to an output port:
 - $PxOUT=x02$
- Reading from an input port:
 - $My_port=P1IN$

Port Registers



Count Characters 'E'

```
;-----  
; File      : Lab4_D1.asm (CPE 325 Lab4 Demo code)  
; Function   : Counts the number of characters E in a given string  
; Description: Program traverses an input array of characters  
;                 to detect a character 'E'; exits when a NULL is detected  
; Input      : The input string is specified in myStr  
; Output     : The port P1OUT displays the number of E's in the string  
; Author     : A. Milenkovic, milenkovic@computer.org  
; Date       : August 14, 2008  
;  
.cdecls C,LIST,"msp430.h"           ; Include device header file  
  
;  
.def    RESET                      ; Export program entry-point to  
; make it known to linker.  
myStr: .string "HELLO WORLD, I AM THE MSP430!", ''  
;  
.text                           ; Assemble into program memory.  
.retain                         ; Override ELF conditional linking  
; and retain current section.  
.retainrefs                     ; And retain any sections that have  
; references to current section.  
;  
RESET:  mov.w  #__STACK_END,SP      ; Initialize stack pointer  
        mov.w  #WDTPW|WDTHOLD,&WDTCTL ; Stop watchdog timer
```

Count Characters 'E' (cont'd)

```
;-----  
; Main loop here  
;  
main:  bis.b  #0FFh,&P1DIR      ; configure P1.x output  
        mov.w  #myStr, R4          ; load the starting address of the string into R4  
        clr.b  R5                ; register R5 will serve as a counter  
  
gnext: mov.b  @R4+, R6           ; get a new character  
        cmp    #0,R6             ; is it a null character  
        jeq   lend              ; if yes, go to the end  
        cmp.b  #'E',R6          ; is it an 'E' character  
        jne   gnext             ; if not, go to the next  
        inc.w  R5                ; if yes, increment counter  
        jmp   gnext             ; go to the next character  
  
lend:   mov.b  R5,&P1OUT      ; set all P1 pins (output)  
        bis.w  #LPM4,SR         ; LPM4  
        nop                           ; required only for Debugger  
  
;  
; Stack Pointer definition  
;  
.global __STACK_END  
.sect   .stack  
;  
; Interrupt Vectors  
;  
        .sect   ".reset"          ; MSP430 RESET Vector  
        .short RESET  
        .end
```

The Case for Subroutines: An Example

- Problem
 - Sum up elements of two integer arrays
 - Display results on P2OUT&P1OUT and P4OUT&P3OUT
- Example
 - arr1 .int 1, 2, 3, 4, 1, 2, 3, 4 ; the first array
 - arr2 .int 1, 1, 1, 1, -1, -1, -1 ; the second array
 - Results
 - P2OUT&P1OUT=0x000A, P4OUT&P3OUT=0x0001
- Approach
 - Input numbers: arrays
 - Main program (no subroutines): initialization, program loops

Sum Up Two Integer Arrays (ver1)

```
;-----  
; File      : Lab5_D1.asm (CPE 325 Lab5 Demo code)  
; Function   : Finds a sum of two integer arrays  
; Description: The program initializes ports,  
;               sums up elements of two integer arrays and  
;               display sums on parallel ports  
; Input      : The input arrays are signed 16-bit integers in arr1 and arr2  
; Output     : P1OUT&P2OUT displays sum of arr1, P3OUT&P4OUT displays sum of arr2  
; Author     : A. Milenkovic, milenkovic@computer.org  
; Date       : September 14, 2008  
;  
.cdecls C,LIST,"msp430.h"           ; Include device header file  
  
;  
.def    RESET                      ; Export program entry-point to  
; make it known to linker.  
;  
.text  
.retain  
.retainrefs  
;  
RESET:   mov.w  #__STACK_END,SP      ; Initialize stack pointer  
StopWDT:  mov.w  #WDTPW|WDTHOLD,&WDTCTL ; Stop watchdog timer
```

Sum up two integer arrays (ver1)

```
;  
; Main code here  
;  
main:    bis.b  #0xFF,&P1DIR          ; configure P1.x as output  
           bis.b  #0xFF,&P2DIR          ; configure P2.x as output  
           bis.b  #0xFF,&P3DIR          ; configure P3.x as output  
           bis.b  #0xFF,&P4DIR          ; configure P4.x as output  
           ; load the starting address of the array1 into the register R4  
           mov.w  #arr1, R4  
           ; load the starting address of the array1 into the register R4  
           mov.w  #arr2, R5  
           ; Sum arr1 and display  
           clr.w  R7  
           mov.w  #8, R10  
           add.w  @R4+, R7  
           dec.w  R10  
           jnz   lnext1  
           mov.b  R7, P1OUT  
           swpb  R7  
           mov.b  R7, P2OUT  
  
lnext1:    ; Holds the sum  
           ; number of elements in arr1  
           ; get next element  
  
           ; display sum of arr1
```

Sum up two integer arrays (ver1)

```
; Sum arr2 and display
    clr.w   R7
    mov.w   #7, R10
    lnext2: add.w   @R5+, R7
    dec.w   R10
    jnz    lnext2
    mov.b   R7, P3OUT
    swpb
    mov.b   R7, P4OUT
    jmp    $

arr1:     .int    1, 2, 3, 4, 1, 2, 3, 4      ; the first array
arr2:     .int    1, 1, 1, 1, -1, -1, -1      ; the second array

; -----
; Stack Pointer definition
; -----
.global __STACK_END
.sect .stack

; -----
; Interrupt Vectors
; -----
.reset
.short RESET
.end
```

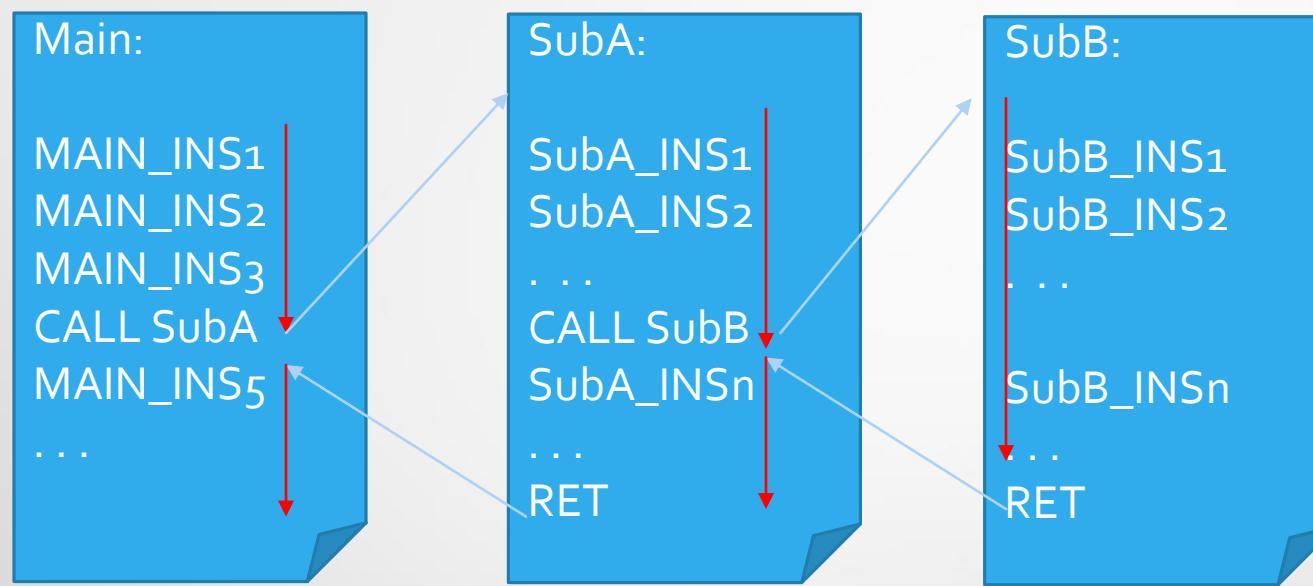
Subroutines

- A particular sub-task is performed many times on different data values
- Frequently used subtasks are known as subroutines
- Subroutines: How do they work?
 - Only one copy of the instructions that constitute the subroutine is placed in memory
 - Any program that requires the use of the subroutine simply branches to its starting location in memory
 - Upon completion of the task in the subroutine, the execution continues at the next instruction in the calling program

Subroutines (cont'd)

- CALL instruction:
perform the branch to subroutines
 - $SP \leq SP - 2$; allocate a word on the stack for return address
 - $M[SP] \leq PC$; push the return address (current PC) onto the stack
 - $PC \leq \text{TargetAddress}$; the starting address of the subroutine is moved into PC
- RET instruction:
the last instruction in the subroutine
 - $PC \leq M[SP]$; pop the return address from the stack
 - $SP \leq SP + 2$; release the stack space

Subroutine Nesting



Mechanisms for Passing Parameters

- Through registers
- Through stack
 - By value
 - Actual parameter is transferred
 - If the parameter is modified by the subroutine, the “new value” does not affect the “old value”
 - By reference
 - The address of the parameter is passed
 - There is only one copy of parameter
 - If parameter is modified, it is modified globally

Subroutine: SUMA_RP

- Subroutine for summing up elements of an integer array
- Passing parameters through registers
 - R12 – starting address of the array
 - R13 – array length
 - R14 – display id
(0 for P2&P1, 1 for P4&P3)

Subroutine: SUMA_RP

```
;-----  
; File      : Lab5_D2_RP.asm (CPE 325 Lab5 Demo code)  
; Function   : Finds a sum of an input integer array  
; Description: suma_rp is a subroutine that sums elements of an integer array  
; Input       : The input parameters are:  
;                 R12 -- array starting address  
;                 R13 -- the number of elements (>= 1)  
;                 R14 -- display ID (0 for P1&P2 and 1 for P3&P4)  
; Output      : No output  
; Author     : A. Milenkovic, milenkovic@computer.org  
; Date       : September 14, 2008  
;-----  
.cdecls C,LIST,"msp430.h"      ; Include device header file  
  
.def suma_rp  
  
.text
```

Subroutine: SUMA_RP

```
suma_rp:  
    push.w  R7          ; save the register R7 on the stack  
    clr.w   R7          ; clear register R7 (keeps the sum)  
lnext:  add.w   @R12+, R7      ; add a new element  
    dec.w   R13         ; decrement step counter  
    jnz    lnext        ; jump if not finished  
    bit.w   #1, R14      ; test display ID  
    jnz    lp34         ; jump on lp34 if display ID=1  
    mov.b   R7, P1OUT     ; display lower 8-bits of the sum on P1OUT  
    swpb    R7          ; swap bytes  
    mov.b   R7, P2OUT     ; display upper 8-bits of the sum on P2OUT  
    jmp    lend         ; skip to end  
lp34:  mov.b   R7, P3OUT     ; display lower 8-bits of the sum on P3OUT  
    swpb    R7          ; swap bytes  
    mov.b   R7, P4OUT     ; display upper 8-bits of the sum on P4OUT  
lend:  pop    R7          ; restore R7  
    ret                 ; return from subroutine  
  
.end
```

Main (ver2): Call suma_rp

```
;-----  
; Main code here  
;  
main:    bis.b    #0xFF,&P1DIR          ; configure P1.x as output  
            bis.b    #0xFF,&P2DIR          ; configure P2.x as output  
            bis.b    #0xFF,&P3DIR          ; configure P3.x as output  
            bis.b    #0xFF,&P4DIR          ; configure P4.x as output  
  
            mov.w    #arr1, R12           ; put address into R12  
            mov.w    #8, R13             ; put array length into R13  
            mov.w    #0, R14             ; display #0 (P1&P2)  
            call     #suma_rp  
  
            mov.w    #arr2, R12           ; put address into R12  
            mov.w    #7, R13             ; put array length into R13  
            mov.w    #1, R14             ; display #0 (P3&P4)  
            call     #suma_rp  
            jmp     $  
  
arr1:    .int     1, 2, 3, 4, 1, 2, 3, 4      ; the first array  
arr2:    .int     1, 1, 1, 1, -1, -1, -1       ; the second array
```

Subroutine: SUMA_SP

- Subroutine for summing up elements of an integer array
- Passing parameters through the stack
 - The calling program prepares input parameters on the stack

Main (ver3): Call suma_sp (Pass Through Stack)

```

; Main code here
;

main:    bis.b   #0xFF,&P1DIR           ; configure P1.x as output
          bis.b   #0xFF,&P2DIR           ; configure P2.x as output
          bis.b   #0xFF,&P3DIR           ; configure P3.x as output
          bis.b   #0xFF,&P4DIR           ; configure P4.x as output

          push    #arr1
          push    #8
          push    #0
          call    #suma_sp
          add.w  #6,SP
          ; collapse the stack

          push    #arr2
          push    #7
          push    #1
          call    #suma_sp
          add.w  #6,SP
          ; collapse the stack

          jmp    $

arr1:    .int    1, 2, 3, 4, 1, 2, 3, 4 ; the first array
arr2:    .int    1, 1, 1, 1, -1, -1, -1 ; the second array

```

Subroutine: SUMA_SP

```
;-----  
; File      : Lab5_D3_SP.asm (CPE 325 Lab5 Demo code)  
; Function   : Finds a sum of an input integer array  
; Description: suma_sp is a subroutine that sums elements of an integer array  
; Input       : The input parameters are on the stack pushed as follows:  
;                  starting address of the array  
;                  array length  
;                  display id  
; Output      : No output  
; Author     : A. Milenkovic, milenkovic@computer.org  
; Date       : September 14, 2008  
;  
.cdecls C,LIST,"msp430.h"           ; Include device header file  
  
.def    suma_sp  
  
.text
```

Subroutine: SUMA_SP (cont'd)

suma_sp:

```

    push   R7
    push   R6
    push   R4
    clr.w  R7
    mov.w  10(SP), R6
    mov.w  12(SP), R4
    lnext:
        add.w @R4+, R7
        dec.w R6
        jnz   lnext
        mov.w 8(SP), R4
        bit.w #1, R4
        jnz   lp34
        mov.b  R7, P1OUT
    P1OUT:
        swpb  R7
        mov.b  R7, P2OUT
        jmp   lend
    lp34:
        mov.b  R7, P3OUT
        swpb  R7
        mov.b  R7, P4OUT
    lend:
        pop   R4
        pop   R6
        pop   R7
        ret
    .end
  
```

; save the registers on the stack
 ; save R7, temporal sum
 ; save R6, array length
 ; save R5, pointer to array
 ; clear R7
 ; retrieve array length
 ; retrieve starting address
 ; add next element
 ; decrement array length
 ; repeat if not done
 ; get id from the stack
 ; test display id
 ; jump to lp34 display id = 1
 ; lower 8 bits of the sum to
 ; swap bytes
 ; upper 8 bits of the sum to P2OUT
 ; jump to lend
 ; lower 8 bits of the sum to P3OUT
 ; swap bytes
 ; upper 8 bits of the sum to P4OUT
 ; restore R4
 ; restore R6
 ; restore R7
 ; return

Address	Stack
0x0800	OTOS
0x07FE	#arr1
0x07FC	0008
0x07FA	0000
0x07F8	Ret. Addr.
0x07F6	(R7)
0x07F4	(R6)
0x07F2	(R4)

The Stack and Local Variables

- Subroutines often need local workspace
- We can use a fixed block of memory space – static allocation – but:
 - The code will not be relocatable
 - The code will not be reentrant
 - The code will not be able to be called recursively
- Better solution: dynamic allocation
 - Allocate all local variables on the stack
 - STACK FRAME = a block of memory allocated by a subroutine to be used for local variables
 - FRAME POINTER = an address register used to point to the stack frame

Subroutine: SUMA_SPSF

```
;-----  
; File      : Lab5_D4_SPSF.asm (CPE 325 Lab5 Demo code)  
; Function   : Finds a sum of an input integer array  
; Description: suma_spsf is a subroutine that sums elements of an integer array.  
;  
;           The subroutine allocates local variables on the stack:  
;  
;           counter (SFP+2)  
;           sum (SFP+4)  
;  
; Input      : The input parameters are on the stack pushed as follows:  
;  
;           starting address of the array  
;  
;           array length  
;  
;           display id  
;  
; Output     : No output  
;  
; Author     : A. Milenkovic, milenkovic@computer.org  
;  
; Date       : September 14, 2008  
;  
;-----  
.cdecls C,LIST,"msp430.h"          ; Include device header file  
  
.def    suma_spsf  
  
.text
```

Subroutine: SUMA_SPSF (cont'd)

```
suma_spsf:  
    ; save the registers on the stack  
    push R12          ; save R12 - R12 is stack frame pointer  
    mov.w SP, R12     ; R12 points on the bottom of the stack frame  
    sub.w #4, SP      ; allocate 4 bytes for local variables  
    push R4          ; pointer register  
    clr.w -4(R12)    ; clear sum, sum=0  
    mov.w 6(R12), -2(R12) ; get array length  
    mov.w 8(R12), R4  ; R4 points to the array starting address  
lnext: add.w @R4+, -4(R12) ; add next element  
    dec.w -2(R12)    ; decrement counter  
    jnz lnext        ; repeat if not done  
    bit.w #1, 4(R12) ; test display id  
    jnz lp34         ; jump to lp34 if display id = 1  
    mov.b -4(R12), P1OUT ; lower 8 bits of the sum to P1OUT  
    mov.b -3(R12), P2OUT ; upper 8 bits of the sum to P2OUT  
    jmp lend         ; skip to lend  
lp34: mov.b -4(R12), P3OUT ; lower 8 bits of the sum to P3OUT  
    mov.b -3(R12), P4OUT ; upper 8 bits of the sum to P4OUT  
lend: pop R4           ; restore R4  
    add.w #4, SP      ; collapse the stack frame  
    pop R12          ; restore stack frame pointer  
    ret              ; return  
.end
```

Address	Stack
0x0800	OTOS
0x07FE	#arr1
0x07FC	0008
0x07FA	0000
0x07F8	Ret. Addr.
0x07F6	(R12)
0x07F4	counter
0x07F2	sum
0x0731	(R4)

Performance

- Performance: how fast a task can be completed
- $\text{Performance}(X) = 1/\text{ExecutionTime}(X)$
- ET: ExecutionTime

$$ET = IC \cdot CPI \cdot CCT = \frac{IC \cdot CPI}{CF}$$

- IC: Instruction Count – the number of instructions executed in the program
- CPI: Cycles Per Instruction – the average number of clock cycles it takes to execute an instruction
- CCT: Clock Cycle Time – the duration of one processor clock cycle
- CF: Clock Frequency (1/CCT)

Performance: An Example

```
RESET:    mov.w   #__STACK_END,SP      ; 4cc
StopWDT:   mov.w   #WDTPW|WDTHOLD,&WDTCTL ; 5cc
            push    R14                ; 3 cc (table 3.15)
            mov.w   SP,R14             ; 1 cc
            mov.w   #aend,R6            ; 2 cc
            mov.w   R6,R5               ; 1 cc
            sub.w   #arr1,R5            ; 2 cc
            sub.w   R5,SP               ; 1 cc
lnext:     dec.w   R6                ; 1 cc x 9
            dec.w   R14               ; 1 cc x 9
            mov.b   @R6,0(R14)          ; 4 cc x 9
            dec.w   R5                ; 1 cc x 9
            jnz    lnext              ; 2 cc x 9
            jmp    $

arr1     .byte   1, 2, 3, 4, 5, 6, 7, 8, 9
aend
.end

TOTAL NUMBER OF CLOCK CYLES:        4+5+3+1+2+1+2+1+9x(1+1+4+1+2) = 19+9x9 = 100 cc
TOTAL NUMBER OF INSTRUCTIONS       8+9x5 = 53 instructions
CPI                                100/53 = 1.88 cc/instruction
```