

# CPE 325: Embedded Systems Laboratory

## Laboratory #5 Tutorial

### MSP430 Assembly Language Programming

#### Subroutines, Passing Parameters, and Hardware Multiplier

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#### Objective:

This tutorial will continue the introduction to assembly language programming with the MSP430 hardware. In this lab, you will learn the following topics:

*Developing subroutines in assembly language*

*Passing parameters to subroutines using registers and the stack*

*Working with hardware multiplier on the MSP430*

#### Notes:

All previous tutorials are required for successful completion of this lab, especially, the tutorials introducing the TI Experimenter's Board and the Code Composer Studio software development environment.

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# 1 Subroutines

In a given program, it is often needed to perform a particular sub-task many times on different data values. Such a subtask is usually called a subroutine. For example, a subroutine may sort numbers in an integer array or perform a complex mathematical operation on an input variable (e.g., calculate  $\sin(x)$ ). It should be noted, that the block of instructions that constitute a subroutine can be included at every point in the main program when that task is needed. However, this would be an unnecessary waste of memory space. Rather, only one copy of the instructions that constitute the subroutine is placed in memory and any program that requires the use of the subroutine simply branches to its starting location in memory. The instruction that performs this branch is named a CALL instruction. The calling program is called CALLER and the subroutine called is called CALLEE.

The instruction that is executed right after the CALL instruction is the first instruction of the subroutine. The last instruction in the subroutine is a RETURN instruction, and we say that the subroutine returns to the program that called it. Since a subroutine can be called from different places in a calling program, we must have a mechanism to return to the appropriate location (the first instruction that follows the CALL instruction in the calling program). At the time of executing the CALL instruction we know the program location of the instruction that follows the CALL (the program counter or PC is pointing to the next instruction). Hence, we should save the return address at the time the CALL instruction is executed. The way in which a machine makes it possible to call and return from subroutines is referred to as its *subroutine linkage method*. The simplest subroutine linkage method is to save the return address in a specific location. This location may be a register dedicated to this function, often referred to as the link register. When the subroutine completes its task, the return instruction returns to the calling program by branching indirectly through the link register.

The CALL instruction is a special branch instruction and performs the following operations:

*Stores the contents of the PC in the link register*

*Branches to the target address specified by the instruction.*

The RETURN instruction is a special branch instruction that performs the following operations:

*Branches to the address contained in the link register.*

## 1.1 Subroutine Nesting

A common programming practice, called subroutine nesting, is to have one subroutine call another. In this case, the return address of the second call is also stored in the link register destroying the previous contents. Hence, it is essential to save the contents of the link register in some other location before calling another subroutine. Subroutine nesting can be carried out to any depth. For example, imagine the following sequence: subroutine A calls subroutine B, subroutine B calls subroutine C, and finally subroutine C calls subroutine D. In this case, the last subroutine D completes its computations and returns to the subroutine C that called it. Next, C completes its execution and returns to the subroutine B that called it and so on. The sequence of returns is as follows: D returns to C, C returns to B, and B returns to A. That is, the return

addresses are generated and used in the last-in-first-out order. This suggests that the return addresses associated with subroutine calls should be pushed onto a stack. Many processors do this automatically. A particular register is designated as the stack pointer, or SP, that is implicitly used in this operation. The stack pointer points to a stack called the processor stack.

The CALL instruction is a special branch instruction and performs the following operations:

- Pushes the contents of the PC on the top of the stack*
- Updates the stack pointer*
- Branches to the target address specified by the instruction*

The RETURN instruction is a special branch instruction that performs the following operations:

- Pops the return address from the top of the stack into the PC*
- Updates the stack pointer.*

## 1.2 Parameter Passing

When calling a subroutine, a calling program needs a mechanism to provide to the subroutine the input parameters, the operands that will be used in computation in the subroutine or their addresses. Later, the subroutine needs a mechanism to return output parameters, the results of the subroutine computation. This exchange of information between a calling program and a subroutine is referred to as parameter passing. Parameter passing may be accomplished in several ways. The parameters can be placed in registers or in memory locations, where they can be accessed by subroutine. Alternatively, the parameters may be placed on a processor stack.

Let us consider the following program shown in Figure 1. We have two integer arrays arr1 and arr2. The program finds the sum of the integers in arr1 and displays the result on the ports P1 and P2, and then finds the sum of the integers in arr2 and displays the result on the ports P3 and P4. It is obvious that we can have a single subroutine that will perform this operation and thus make our code more readable and reusable. The subroutine needs to get three input parameters: what is the starting address of the input array, how many elements the array has, and where to display the result. In this example, the subroutine does not return any output parameter to the calling program.

```
1 ;-----
2 ; File      : Lab5_D1.asm (CPE 325 Lab5 Demo code)
3 ; Function   : Finds a sum of two integer arrays
4 ; Description: The program initializes ports,
5 ;             sums up elements of two integer arrays and
6 ;             display sums on parallel ports
7 ; Input      : The input arrays are signed 16-bit integers in arr1 and arr2
8 ; Output     : P1OUT&P2OUT displays sum of arr1, P3OUT&P4OUT displays sum of arr2
9 ; Author     : A. Milenkovic, milenkovic@computer.org
10 ; Date      : September 14, 2008
11 ;-----
12             .cdecls C,LIST,"msp430.h"           ; Include device header file
13
14 ;-----
15             .def      RESET                       ; Export program entry-point to
16                                     ; make it known to linker.
17 ;-----
```

```

18      .text                ; Assemble into program memory.
19      .retain              ; Override ELF conditional linking
20                                ; and retain current section.
21      .retainrefs         ; And retain any sections that have
22                                ; references to current section.
23
24 ;-----
25 RESET:    mov.w    #__STACK_END,SP    ; Initialize stack pointer
26 StopWDT:  mov.w    #WDTPW|WDTHOLD,&WDTCTL ; Stop watchdog timer
27
28 ;-----
29 ; Main code here
30 ;-----
31 main:    bis.b    #0xFF, &P1DIR        ; configure P1.x as output
32            bis.b    #0xFF, &P2DIR        ; configure P2.x as output
33            bis.b    #0xFF, &P3DIR        ; configure P3.x as output
34            bis.b    #0xFF, &P4DIR        ; configure P4.x as output
35            ; load the starting address of the array1 into the register R4
36            mov.w    #arr1, R4
37            ; load the starting address of the array2 into the register R5
38            mov.w    #arr2, R5
39            ; Sum arr1 and display
40            clr.w    R7                    ; holds the sum
41            mov.w    #8, R10                ; number of elements in arr1
42 lnext1:  add.w    @R4+, R7                ; add the current element to sum
43            dec.w    R10                    ; decrement arr1 length
44            jnz     lnext1                  ; get next element
45            mov.b    R7, P1OUT              ; display lower byte of sum of arr1
46            swpb    R7                    ; swap bytes
47            mov.b    R7, P2OUT              ; display upper byte of sum of arr1
48            ; Sum arr2 and display
49            clr.w    R7                    ; Holds the sum
50            mov.w    #7, R10                ; number of elements in arr2
51 lnext2:  add.w    @R5+, R7                ; get next element
52            dec.w    R10                    ; decrement arr2 length
53            jnz     lnext2                  ; get next element
54            mov.b    R7, P3OUT              ; display lower byte of sum of arr2
55            swpb    R7                    ; swap bytes
56            mov.b    R7, P4OUT              ; display upper byte of sum of arr2
57            jmp     $
58
59 arr1:    .int    1, 2, 3, 4, 1, 2, 3, 4    ; the first array
60 arr2:    .int    1, 1, 1, 1, -1, -1, -1    ; the second array
61
62 ;-----
63 ; Stack Pointer definition
64 ;-----
65            .global __STACK_END
66            .sect   .stack
67
68 ;-----
69 ; Interrupt Vectors
70 ;-----
71            .sect   ".reset"                ; MSP430 RESET Vector
72            .short  RESET

```

**Figure 1. Array Addition without a Subroutine (Lab5\_D1.asm)**

Let us next consider the main program (Figure 2) where we pass the parameters through registers. Passing parameters through the registers is straightforward and efficient. Three input parameters are placed in registers as follows: R12 keeps the starting address of the input array, R13 keeps the array length, and R14 defines the display identification (#0 for P1&P2 and #1 for P3&P4). The calling program places the parameters in these registers, and then calls the subroutine using the CALL #suma\_rp instruction. The subroutine shown in Figure 3 uses register R7 to hold the sum of the integers in the array. The register R7 may contain valid data that belongs to the calling program, so our first step should be to push the content of the register R7 on the stack. The last instruction before the return from the subroutine is to restore the original content of R7. Generally, it is a good practice to save all the general-purpose registers used as temporary storage in the subroutine as the first thing in the subroutine, and to restore their original contents (the contents pushed on the stack at the beginning of the subroutine) just before returning from the subroutine. This way, the calling program will find the original contents of the registers as they were before the CALL instruction. Other registers that our subroutine uses are R12, R13, and R14. These registers keep parameters, so we assume we can modify them (they do not need to preserve their original value once we are back in the calling program).

```

1 ;-----
2 ; File      : Lab5_D2_main.asm (CPE 325 Lab5 Demo code)
3 ; Function  : Finds a sum of two integer arrays using subroutines
4 ; Description: The program initializes ports and
5 ;           : calls suma_rp to sum up elements of integer arrays and
6 ;           : display sums on parallel ports.
7 ;           : Parameters to suma_rp are passed through registers, R12, R13, R14.
8 ; Input    : The input arrays are signed 16-bit integers in arr1 and arr2
9 ; Output   : P1OUT&P2OUT displays sum of arr1, P3OUT&P4OUT displays sum of arr2
10 ; Author   : A. Milenkovic, milenkovic@computer.org
11 ; Date    : September 14, 2008
12 ;-----
13         .cdecls C,LIST,"msp430.h"           ; Include device header file
14
15 ;-----
16         .def      RESET                     ; Export program entry-point to
17                                     ; make it known to linker.
18         .ref      suma_rp
19 ;-----
20         .text                               ; Assemble into program memory.
21         .retain                               ; Override ELF conditional linking
22                                     ; and retain current section.
23         .retainrefs                          ; And retain any sections that have
24                                     ; references to current section.
25 ;-----
26 RESET:      mov.w   #__STACK_END,SP        ; Initialize stack pointer
27 StopWDT:   mov.w   #WDTPW|WDTHOLD,&WDTCTL ; Stop watchdog timer
28
29 ;-----

```

```

30 ; Main code here
31 ;-----
32 main:    bis.b    #0xFF,&P1DIR    ; configure P1.x as output
33         bis.b    #0xFF,&P2DIR    ; configure P2.x as output
34         bis.b    #0xFF,&P3DIR    ; configure P3.x as output
35         bis.b    #0xFF,&P4DIR    ; configure P4.x as output
36
37         mov.w    #arr1, R12      ; put address into R12
38         mov.w    #8, R13        ; put array length into R13
39         mov.w    #0, R14        ; display #0 (P1&P2)
40         call    #suma_rp
41
42         mov.w    #arr2, R12      ; put address into R12
43         mov.w    #7, R13        ; put array length into R13
44         mov.w    #1, R14        ; display #0 (P3&P4)
45         call    #suma_rp
46         jmp     $
47
48 arr1:    .int     1, 2, 3, 4, 1, 2, 3, 4 ; the first array
49 arr2:    .int     1, 1, 1, 1, -1, -1, -1 ; the second array
50
51 ;-----
52 ; Stack Pointer definition
53 ;-----
54         .global  __STACK_END
55         .sect   .stack
56
57 ;-----
58 ; Interrupt Vectors
59 ;-----
60         .sect   ".reset"        ; MSP430 RESET Vector
61         .short  RESET
62         .end

```

Figure 2. Array Addition Using `suma_rp` Subroutine (Lab5\_D2\_main.asm)

```

1 ;-----
2 ; File      : Lab5_D2_RP.asm (CPE 325 Lab5 Demo code)
3 ; Function  : Finds a sum of an input integer array
4 ; Description: suma_rp is a subroutine that sums elements of an integer array
5 ; Input    : The input parameters are:
6 ;           R12 -- array starting address
7 ;           R13 -- the number of elements (>= 1)
8 ;           R14 -- display ID (0 for P1&P2 and 1 for P3&P4)
9 ; Output   : No output
10 ; Author   : A. Milenkovic, milenkovic@computer.org
11 ; Date    : September 14, 2008
12 ;-----
13         .cdecls C,LIST,"msp430.h" ; Include device header file
14
15         .def suma_rp
16
17         .text
18

```

```

19  suma_rp:
20      push.w  R7          ; save the register R7 on the stack
21      clr.w   R7          ; clear register R7 (keeps the sum)
22  lnext:  add.w   @R12+, R7 ; add a new element
23      dec.w   R13         ; decrement step counter
24      jnz    lnext       ; jump if not finished
25      bit.w   #1, R14     ; test display ID
26      jnz    lp34        ; jump on lp34 if display ID=1
27      mov.b  R7, P1OUT    ; display lower 8-bits of the sum on P1OUT
28      swpb   R7          ; swap bytes
29      mov.b  R7, P2OUT    ; display upper 8-bits of the sum on P2OUT
30      jmp    lend        ; skip to end
31  lp34:  mov.b  R7, P3OUT  ; display lower 8-bits of the sum on P3OUT
32      swpb   R7          ; swap bytes
33      mov.b  R7, P4OUT    ; display upper 8-bits of the sum on P4OUT
34  lend:  pop    R7        ; restore R7
35      ret                    ; return from subroutine
36      .end

```

**Figure 3. Subroutine that Adds up the Elements of the Array (Lab5\_D2\_RP.asm)**

If many parameters are passed, there may not be enough general-purpose registers available for passing parameters into the subroutine. In this case we use the stack to pass parameters. Figure 4 shows the calling program (Lab5\_D3\_main.asm) and Figure 5 shows the subroutine (Lab5\_D3\_SP.asm). Before calling the subroutine, we place parameters on the stack using PUSH instructions (the array starting address, array length, and display id – each parameter is 2 bytes long). The CALL instruction pushes the return address on the stack. The subroutine then stores the contents of the registers R7, R6, and R4 on the stack (another 8 bytes) to save their original content. The next step is to retrieve input parameters (array starting address and array length). They are on the stack, but to know exactly where, we need to know the current state of the stack and its organization (how does it grow, and where does SP point to). The original values of the registers pushed onto the stack occupy 6 bytes, the return address 2 bytes, the display id 2 bytes, and the array length 2 bytes. The total distance between the top of the stack and the location on the stack where we placed the starting address is 12 bytes. So the instruction MOV 12(SP), R4 loads the register R4 with the first parameter (the array starting address). Similarly, the array length can be retrieved by MOV 10(SP), R6. The register values are restored before returning from the subroutine (notice the reverse order of POP instructions). Once we are back in the calling program, we can free 6 bytes on the stack used to pass parameters.

```

1  ;-----
2  ; File       : Lab5_D3_main.asm (CPE 325 Lab5 Demo code)
3  ; Function   : Finds a sum of two integer arrays using a subroutine suma_sp
4  ; Description: The program initializes ports and
5  ;             calls suma_rp to sum up elements of integer arrays and
6  ;             display sums on parallel ports.
7  ;           Parameters to suma_sp are passed through the stack.
8  ; Input     : The input arrays are signed 16-bit integers in arr1 and arr2
9  ; Output    : P1OUT&P2OUT displays sum of arr1, P3OUT&P4OUT displays sum of arr2
10 ; Author    : A. Milenkovic, milenkovic@computer.org
11 ; Date     : September 14, 2008
12 ;-----

```

```

13         .cdecls C,LIST,"msp430.h"           ; Include device header file
14
15 ;-----
16         .def      RESET                       ; Export program entry-point to
17                                         ; make it known to linker.
18         .ref      suma_sp
19 ;-----
20         .text                                   ; Assemble into program memory.
21         .retain                                ; Override ELF conditional linking
22                                         ; and retain current section.
23         .retainrefs                            ; And retain any sections that have
24                                         ; references to current section.
25 ;-----
26 RESET:    mov.w   #__STACK_END,SP           ; Initialize stack pointer
27 StopWDT: mov.w   #WDTPW|WDTHOLD,&WDTCTL    ; Stop watchdog timer
28
29 ;-----
30 ; Main code here
31 ;-----
32 main:    bis.b   #0xFF,&P1DIR                ; configure P1.x as output
33         bis.b   #0xFF,&P2DIR                ; configure P2.x as output
34         bis.b   #0xFF,&P3DIR                ; configure P3.x as output
35         bis.b   #0xFF,&P4DIR                ; configure P4.x as output
36
37         push    #arr1                        ; push the address of arr1
38         push    #8                          ; push the number of elements
39         push    #0                          ; push display id
40         call    #suma_sp
41         add.w   #6,SP                        ; collapse the stack
42         push    #arr2                        ; push the address of arr1
43         push    #7                          ; push the number of elements
44         push    #1                          ; push display id
45         call    #suma_sp
46         add.w   #6,SP                        ; collapse the stack
47
48         jmp     $
49
50 arr1:    .int     1, 2, 3, 4, 1, 2, 3, 4      ; the first array
51 arr2:    .int     1, 1, 1, 1, -1, -1, -1     ; the second array
52
53 ;-----
54 ; Stack Pointer definition
55 ;-----
56         .global  __STACK_END
57         .sect   .stack
58
59 ;-----
60 ; Interrupt Vectors
61 ;-----
62         .sect   ".reset"                     ; MSP430 RESET Vector
63         .short  RESET
64         .end

```

**Figure 4. Example of Passing Parameters to Subroutine Through the Stack (Lab5\_D3\_main.asm)**



```

1 ;-----
2 ; File      : Lab5_D3_SP.asm (CPE 325 Lab5 Demo code)
3 ; Function   : Finds a sum of an input integer array
4 ; Description: suma_sp is a subroutine that sums elements of an integer array
5 ; Input      : The input parameters are on the stack pushed as follows:
6 ;             starting address of the array
7 ;             array length
8 ;             display id
9 ; Output     : No output
10 ; Author    : A. Milenkovic, milenkovic@computer.org
11 ; Date      : September 14, 2008
12 ;-----
13             .cdecls C,LIST,"msp430.h"          ; Include device header file
14
15             .def      suma_sp
16
17             .text
18 suma_sp:
19
20             push     R7          ; save the registers on the stack
21             push     R6          ; save R7, temporal sum
22             push     R4          ; save R6, array length
23             clr.w    R7          ; save R5, pointer to array
24             mov.w    10(SP), R7  ; clear R7
25             mov.w    12(SP), R4  ; retrieve array length
26 lnext:      add.w    @R4+, R7    ; retrieve starting address
27             dec.w    R6          ; add next element
28             jnz     lnext       ; decrement array length
29             mov.w    8(SP), R4   ; repeat if not done
30             bit.w    #1, R4      ; get id from the stack
31             jnz     lp34        ; test display id
32             mov.b    R7, P10UT   ; jump to lp34 display id = 1
33             swpb     R7          ; lower 8 bits of the sum to P10UT
34             mov.b    R7, P20UT   ; swap bytes
35             jmp     lend         ; upper 8 bits of the sum to P20UT
36 lp34:      mov.b    R7, P30UT   ; jump to lend
37             swpb     R7          ; lower 8 bits of ths sum to P30UT
38             mov.b    R7, P40UT   ; swap bytes
39 lend:      mov.b    R7, P40UT   ; upper 8 bits of the sum to P40UT
40             pop     R4          ; restore R4
41             pop     R6          ; restore R6
42             pop     R7          ; restore R7
43             ret          ; return
             .end

```

Figure 5. Array Addition Subroutine that Uses Parameters from the Stack (Lab5\_D3\_SP.asm)

## 2 Hardware Multiplier

The MSP430 contains an optional peripheral hardware multiplier that allows the user to quickly perform multiplication operations. Multiplication operations using the standard instruction set

can be complex and consume a lot of processing time; however, the hardware multiplier is a specialized peripheral that the user can operate with only a few instructions. The multiplier can perform up to 16-bit by 16-bit multiplication and can perform signed or unsigned multiplication with or without an accumulator. Some MSP430 models have no multiplier, but some models have a 32-bit by 32-bit multiplier. It is important to check the datasheet for your particular device to understand the available peripherals.

To use the hardware multiplier, you simply move your first operand (multiplicand) into a register designed to accept the first operand. There are four registers which can accept the first operand, and the one you choose determines the type of multiplication that will be performed. The second operand is then moved to the OP2 register. The result of the multiplication is calculated and placed in two registers – RESLO and RESHI. An additional result register, SUMEXT, is used in certain multiplication operations. The MSP430 user's guide includes a list of examples for performing the different types of multiplication, and they are listed here for convenience.

```

; 16x16 Unsigned Multiply
MOV    #01234h,&MPY    ; Load first operand
MOV    #05678h,&OP2    ; Load second operand
; ...                ; Process results

; 8x8 Unsigned Multiply. Absolute addressing.
MOV    #012h,&0130h    ; Load first operand
MOV    #034h,&0138h    ; Load 2nd operand
; ...                ; Process results

; 16x16 Signed Multiply
MOV    #01234h,&MPYS   ; Load first operand
MOV    #05678h,&OP2   ; Load 2nd operand
; ...                ; Process results

; 8x8 Signed Multiply. Absolute addressing.
MOV.B  #012h,&0132h    ; Load first operand
SXT    &MPYS           ; Sign extend first operand
MOV.   B #034h,&0138h  ; Load 2nd operand
SXT    &OP2            ; Sign extend 2nd operand
; (triggers 2nd multiplication)
; ...                ; Process results

; 16x16 Unsigned Multiply Accumulate
MOV    #01234h,&MAC    ; Load first operand
MOV    #05678h,&OP2    ; Load 2nd operand
; ...                ; Process results

; 8x8 Unsigned Multiply Accumulate. Absolute addressing
MOV.B  #012h,&0134h    ; Load first operand
MOV.B  #034h,&0138h    ; Load 2nd operand
; ...                ; Process results

; 16x16 Signed Multiply Accumulate
MOV    #01234h,&MACS   ; Load first operand
MOV    #05678h,&OP2   ; Load 2nd operand

```

```

; ...                ; Process results

; 8x8 Signed Multiply Accumulate. Absolute addressing
MOV.B #012h,&0136h    ; Load first operand
SXT    &MACS          ; Sign extend first operand
MOV.B #034h,R5        ; Temp. location for 2nd operand
SXT    R5             ; Sign extend 2nd operand
MOV    R5,&OP2        ; Load 2nd operand
; ...                ; Process results

```

### 3 References

You should read the following references to gain more familiarity with subroutines, passing parameters, and the hardware multiplier:

- [MSP430 Assembly Language Programming](#)
- Page 177-185 in Davies' *MSP430 Microcontroller Basics* (subroutines and passing parameters)
- Chapter 8, pages 345-352, in the MSP430FG4618 user's guide (16-bit hardware multiplier)

